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AP/ITEC 2210 3.0 A: System Administration Fall 2023

Instructor: Jamon Camisso ITEC 2210 Chat: mattermost.itec2210.ca Email: jamon@vorku.ca Website: https://eclass.vorku.ca/

Date/Time: Wednesday, 19:00-22:00 Location: Zoom / ACE 003 Office hours: Via Mattermost any time

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– Final exam:

Date - location TBD

• In person

• Open book

• Cumulative

- Backups review:

• The importance of offsite backups:

https://www.reuters.com/article/us-france-ovh-fire/fire-breaks -out-in-ovh-building-in-strasbourg-france-idUSKBN2B20NU

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More Information

First some notes:

 I *will* be oversimplifying or eliding details of some algorithms or maths in this Class - good crypto is hard!

 However, as an SA the principles are important, regardless of implementation details, and I'm confident explaining those, so here goes!

 The nature of cryptography (and security) is it has to be nearly perfect, or it will eventually fail to a smarter or better financed adversary

- Cryptography principle 0
 - Never design your own cryptography system
 - You will get it wrong, someone will find a weakness
 - Idea is affectionately called **Schneier's Law**
 - An example analogy:

 If you wouldn't design and strap yourself to your own rocket, you probably shouldn't design your own cryptosystem

- Cryptography resources
 Menezes, A. J., C., V. O., & Vanstone, S. A. (2001). Handbook of applied cryptography. Boca Raton: CRC Press.
 Available free: <u>https://cacr.uwaterloo.ca/hac/</u>
 - <u>https://crypto.stackexchange.com</u> It is the best kind of pedantry The links to primary and secondary sources are especially helpful
 - NIST Federal Information Processing Standards (FIPS)
 - Wiki.openssl.org
 - To prove crypto is 'fun' see <u>Schneier Facts</u>

Cryptography resources

 Full course in breaking crypto systems (in order to learn about them in a practical way) here: <u>https://cryptopals.com/</u>

Cryptographic attacks

I'm not going to cover them since there are so many, but here are some side-channel (non-cryptographic) attacks that can be more effective than finding mathematical weaknesses in a crypto system Black-bag cryptanalysis - physical theft Man-in-the-middle attack - eavesdropping

Power analysis - what it says, crazy cool

<u>Replay attack - replay encrypted data, e.g. WEP, WPA Handshake</u> <u>Rubber-hose cryptanalysis - physical or psychological torture</u> Timing analysis - using CPU timings to glean information

- Cryptography basics

Alice, Bob, Eve are the main actors who we'll encounter
 Typical scenario is Alice and Bob want to communicate privately, and Eve wants to know what they're saying

- **Plaintext** (message) is what they're communicating
- **Ciphertext** is the encrypted version of the message

Cryptography basics

 Key is the secret that is mixed with plaintext in an encryption algorithm to produce a ciphertext

Keyspace is the size of all possible key values, e.g. 2²⁰⁴⁸

 Key schedule is an algorithm that expands a key from say, 128 bits to thousands of bits for later use in an encryption operation

- Block cipher algorithms operate on chunks of data at a time
- Stream cipher algorithms operate on bits of data at a time

What does encryption Look like?

From HAC ch1 p.13

E - encryption function D - decryption function m - message c - ciphertext



Figure 1.6: Schematic of a two-party communication using encryption.

- Cryptographic **operations**:

Substitution: A=C, B=D, C=E, D=F - Caesar cipher

Transposition: plaintext is reordered

• Bitwise operation: bits are XOR'ed

• **Round**: composition of operation(s) on a plaintext

- Cryptographic **operations** (from HAC, ch1, p.20):

- Substitution in a round adds confusion to a ciphertext
 - Idea is to make the relationship between encryption key and ciphertext appear complex
- **Transposition** adds **diffusion** to a ciphertext
 - Idea is to rearrange and spread out non-uniform bits in a message to statistical patterns, common letter combinations

Crail Division all ranction bla.c.d) Lif(d bla.e.d)

- Cryptography basics

1 XOR 1 = 0
1 XOR 0 = 1
0 XOR 1 = 1
0 XOR 0 = 0

 Tie in to last week's RAID5/6 Class and how XOR & parity works: https://igoro.com/archive/how-raid-6-dual-parity-calculation-works

- Cryptographic Hash functions
- Map an input of any arbitrary length binary string to a fixed length unique binary string (hash) that represents the input
- Take the string `itec2210.ca` and run it through a hashing algorithm (use a terminal on your VM):
 - echo 'itec2210.ca' |md5
 b44f302d2e8f77c54455ce7b7b6f08ba
 - python -c 'print("A"*100)' |md5
 0d6aa8e34b2af058d6e4a27d7ff511dd

- Cryptographic Hash functions
- python -c 'print("A"*100)' |md5
 0d6aa8e34b2af058d6e4a27d7ff511dd
- No matter how long the input, output is always 128 bits
- Likewise with SHA1/2/3, e.g:
- python -c 'print("A"*100)' | shasum -a 256
 f76262cff1f14a9eb92d57ac0de2b8dc431f3c0fe3e88f9394140f6802406
 6b7 -

- Cryptographic Hash functions
- Now repeat the same input to the algorithm:
- python -c 'print("A"*100)' | shasum -a 256
 f76262cff1f14a9eb92d57ac0de2b8dc431f3c0fe3e88f9394140f6802406
 6b7 -
- python -c 'print("A"*100)' | shasum -a 256
 f76262cff1f14a9eb92d57ac0de2b8dc431f3c0fe3e88f9394140f6802406
 6b7 -
- python -c 'print("A"*100)' | shasum -a 256
 f76262cff1f14a9eb92d57ac0de2b8dc431f3c0fe3e88f9394140f6802406
 6b7 -

- Cryptographic Hash functions
- Notice how the output never changes?
- The same input will always produce the same set of output bits
- Enables checking things like:
 - File contents are unmodified
 - Encrypted data is intact
 - Email message signatures
 - Password entries in databases
 - TLS browser encryption

- Cryptographic Hash functions
- Map an input of any arbitrary length binary string to a fixed length unique binary string (hash) that represents the input
- It should be difficult to find two separate inputs that map to the same hash output (a collision) (SHA-1 bruteforce odds being 2^{^80}, or 1.2x10^{^24})
 1,200,000,000,000,000,000,000
- Moreover, given a hash, it should be infeasible to work backwards and find an input: algorithm is not reversible
- But wait:

Cryptographic Hash functions

Common hashing functions you'll encounter:

- o MD5 <u>BROKEN</u>
- SHA-1 BROKEN
- <u>SHA-2</u> current standard but getting weaker against attacks
- **SHA-3** recently released (as of 2015), variable length digests
- **Bcrypt** can be made to run arbitrarily slow to prevent bruteforce
- Hash functions have lifetimes of decades at best
- As computer power grows, speed of finding collisions does too

- Cryptographic Hash functions
- Used extensively to create digital signatures, like SSL/TLS certificates, document signatures, email encryption
- Also used to ensure data integrity, like file transfers over untrusted networks, e.g. Debian, Ubuntu, RedHat, Fedora, software repositories
 - http://centos.mirror.iweb.ca/7/extras/x86_64/repodata/repomd.xml.asc
 - http://centos.mirror.iweb.ca/7/extras/x86_64/repodata/repomd.xml
- Note: signature is a hash of the data, which is then encrypted with the private key - later we'll see how the public key decrypts the signature

- Symmetric and Asymmetric encryption
- Symmetric both parties share a secret encryption key
 Can be block or stream based
 - AES is the current generally accepted standard for symmetric (block) encryption
 - **AES** encryption algorithm uses 16 byte blocks
 - Example would be a passphrase on an encrypted disk volume
 - RC4 is an example of a (broken) stream cipher
 WEP encryption used RC4 in Initialization Vector (IV)
 - Many block ciphers can be used in stream mode

- Symmetric and Asymmetric encryption
- <u>IV aside</u> visit wikipedia article, it is very helpful

 Initialization vector or starting value is used to ensure that repeated use of a key on duplicate plaintext results in different ciphertext

- Usually IV is mixed with key in the algorithm's state machine as part of a key schedule or keystream
- In most block cipher modes IV will be stored with a block of ciphertext

internet in the internet internet

- Symmetric encryption

his[a]}, pushStack: function[a,b

else if(c.attachEvent)(C.a

– HAC ch1 p.16



Symmetric Encryption - AES

- Designed by extensive consultation with cryptography community and standards bodies
- Intent was to replace DES (2⁵⁶ bit key combinations, 7.2x10¹⁶)

– AES has:

- **2^128** or 3.4 x 10³⁸ possible keys (128-bit);
- **2^192** or 6.2×10^{57} possible keys (192-bit); and
- **2^256** or 1.1 x 10⁷⁷ possible keys (256-bit)



 "Assuming that one could build a machine that could recover a DES key in a second (i.e., try 2⁵⁵ keys per second)

 then it would take that machine approximately 149 thousand-billion (149 trillion) years to crack a 128-bit AES key.

 To put that into perspective, the universe is believed to be less than 20 billion years old" (Since AES was published, <u>WMAP</u> <u>determined it is likely 13.77 ± 0.059 billion years</u>)

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 - AES
 - So what does AES actually look like?
 - <u>https://formaestudio.com/rijndaelinspector/archivos/Rijndael_Animation_v4_eng-html5.html</u>
 - It is worth taking your time with it

- Symmetric and Asymmetric encryption
- Asymmetric (AKA Public Key)
 - Designed around a pair of large prime numbers (key pair) and some modular arithmetic, with one number being private, and the other public
 - The public key is used to encrypt a value, which only the private key can decrypt
 - Example would be **RSA** algorithm, used for **PGP/GPG**
 - Think of it like a locked mailbox where anyone can put a letter in, but can't take a letter out without the key

Class 7 - Encryption, Security using the second s

- Recall symmetric encryption

– HAC ch1 p.16

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is[a]}.pushStack:function[s.b.c]

This is asymmetric encryption

– HAC ch1 p.28





Send the message to the recipient..

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Cryptographic hash used for digital signature

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https://www.ibm.com/support/knowledgecenter/SSFKSJ_7.1.0/com.ibm.mg.doc/sy10520_.htm

Asymmetric encryption

Public key can encrypt a message that can only be decrypted with the private key

Used for secure communication

 Private key can encrypt a message that can only be decrypted with the public key
 Used for authentication and signatures

- Bear with me here, the idea is to see that the following works
- Recall modulus arithmetic: $10 \mod 3 \equiv 1$

- Asymmetric RSA encryption (from HAC ch8 p.286)
- Algorithm components are:
 - 1. Generate two large random (and distinct) primes p and q, each roughly the same size
 - 2. Compute n = pq and $\varphi = (p 1)(q 1)$
 - 3. Select a random integer e, $1 < e < \phi$, such that $gcd(e, \phi) = 1$ (e is coprime with phi, that is, their only common divisor is 1)
 - 4. Use the extended Euclidean algorithm (Algorithm 2.107) to compute the unique integer d, 1 < d < φ, such that ed ≡ 1 (mod φ)
 Otherwise expressed as ed / φ = 1
 - 5. A's public key is (n, e); A's private key is d

- Asymmetric RSA encryption (from HAC ch8 p.286)

– Example:

<u>https://simple.wikipedia.org/wiki/RSA_(algorithm)</u>

- Asymmetric RSA encryption (from HAC ch8)

- That e part

3 is a good candidate because it makes small values for computation, but can lead to broken encryption if the same message is sent to multiple recipients
 65537 is a good value, because it is 0x10001 in hex, or 100000000000001 in binary which makes for easy computation
 Take a look at an OpenSSL generated RSA certificate, it is likely that it will use this value. Also 'openssl genrsa -h'
 Many RSA implementations choose e first, then check that phi is coprime with e

Asymmetric RSA encryption

- Alice and Bob can now communicate securely with each other by encrypting messages with the other person's public key
- They can also sign their own messages to prove to the other person that they were the original sender, and that the encrypted message was not tampered with in transit
- Eve can listen in all she likes, and if Bob and Alice chose some large primes, Eve probably won't decrypt any messages until after the heat death of the universe

– Now, brief X.509/SSL/TLS certificate aside - remember this?



om.com/support/knowledgecenter/SSFKSJ_7.1.0/com.ibm.mg.doc/sy10520_.htm

- Now, brief X.509/SSL/TLS certificate aside
- RSA is one possible algorithm in this signing and hashing scheme
- 'Plaintext' is set of fields, like domain name, location, expiry dates
- As well as e and n from the RSA keypair (public key)
- A 3rd party Certificate Authority (CA) digitally signs the certificate with their private key, essentially validating you are who you say you are in the certificate

- Now, brief X.509/SSL/TLS certificate aside
- /etc/ssl/certs on your VM contains 3rd party CA (certificate authority)
 certs, tools like curl, apt rely on these for security
- Your browser and OS will also contain a vetted list of CA certificates
- You can validate a certificate that a CA has signed with their private RSA key based on the principle of reversing encryption using the public key
- Say, for example, a bank's TLS certificate, which itself has a public RSA key in it. Use the key to decrypt signatures, and encrypt messages

- Asymmetric encryption
- Ok, back to Bob and Alice
- What if Alice and Bob have never met, but want to exchange messages using a shared secret key?

– Enter Diffie-Hellman algorithm



https://en.wikipedia.org/wiki/Diffie%E2%80%93Hellman_key_exchange

– Diffie-Hellman key exchange

– <u>Logjam vulnerability</u>

- <u>See also cloudflare</u>



- Putting it all together

 What if two parties who have never met would like to exchange secure messages? Say a credit card transaction online

Recall our tools so far:

- Hash functions SHA-1, SHA-2, SHA-3, MD5
- Symmetric ciphers block and stream, AES, ChaCha20
- Asymmetric encryption
 - RSA algorithm
 - DH algorithm
- X509/SSL/TLS certificates containing CA signed public keys

<u>TLS handshake</u>

"When a TLS client and server first start communicating, they agree on a protocol version, select cryptographic algorithms, optionally authenticate each other, and use public-key encryption techniques to generate shared secrets."



Class 7 - Encryption, Security level:m.level}}};or(j=0,k-p.le

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SSL Handshake (Diffie-Hellman) Without Keyless SSL Handshake



Best Explanation I've seen:

o <u>https://tls12.xargs.org/</u>

– To be continued next week!