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AP/ITEC 2210 3.0 A: System Administration Fall 2023

Instructor: Jamon Camisso ITEC 2210 Chat: mattermost.itec2210.ca Email: jamon@vorku.ca Website: https://eclass.vorku.ca/

Date/Time: Wednesday, 19:00-22:00 Location: Zoom / ACE 003 Office hours: Via Mattermost any time

- Encryption Review
 Symmetric (AES)
 - **Asymmetric**/Public (RSA, Diffie-Hellman)
 - Hash Algorithms (MD5, SHA-{1,2,3})
 - Signatures Private key encrypts a message hash, public key decrypts it, ensuring message authenticity and sender identity, no tampering during transit

- Encryption Review
 Certificate Authorities (CA) act as a trusted third party
 - CA signs your public key with their private key
 - You can then issue signatures, or receive encrypted messages, for example a TLS webserver
 - If a user has the CA's public certificate, they can reasonably trust your certificate

- Encryption Review
 - An alternate model to third parties: <u>Web of Trust</u>
 - A public key collects signatures from various sources
 - With enough independent signatures you can be reasonably sure the public key belongs to someone
 - However, not used in browsers for TLS (mostly email)

– RSA - James Grime

Good review

Practical sized #s



- Now, brief X.509/SSL/TLS certificate aside
- RSA is used in this signing and hashing scheme*
- 'Plaintext' is set of fields, like domain name, location, expiry dates
- As well as e and n from the RSA keypair (public key)
- A 3rd party Certificate Authority (CA) digitally signs the certificate with their private key, essentially validating you are who you say you are in the certificate
- * Elliptic curve and other algorithms can be used, not as widely supported

- Now, brief X.509/SSL/TLS certificate aside
- /etc/ssl/certs on your VM contains 3rd party CA (certificate authority)
 certs, tools like curl, apt rely on these for security
- Your browser and OS will also contain a vetted list of CA certificates
- You can validate a certificate that a CA has signed with their private RSA key based on the principle of reversing encryption using the public key
- Say, for example, a bank's TLS certificate, which itself has a public RSA key in it. Use the key to decrypt signatures, and encrypt messages

- Asymmetric encryption
- Ok, back to Bob and Alice
- What if Alice and Bob have never met, but want to exchange messages using a shared secret key?

a,b){return.e.each(this.a.b)}, ready

Asymmetric review





- Putting it all together

 What if two parties who have never met would like to exchange secure messages? Say a credit card transaction online

Recall our tools so far:

- Hash functions SHA-1, SHA-2, SHA-3, MD5
- Symmetric ciphers block and stream, AES, ChaCha20
- Asymmetric encryption
 - RSA algorithm
 - DH algorithm
- X509/SSL/TLS certificates containing CA signed public keys

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– TLS 1.2 handshake

- <u>RFC 5246 Section-7.3</u>

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Client		Server
ClientHello	>	
		ServerHello
		Certificate [*]
		ServerKeyExchange ¹
		CertificateRequest ¹
	<	ServerHelloDon
Certificate*		
ClientKeyExchange		
CertificateVerify*		
[ChangeCipherSpec]		
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Figure 1. Message flow for a full handshake

– TLS 1.2 handshake

Exchange hello messages to agree on algorithms, exchange random values, and check for session resumption

 Exchange the necessary cryptographic parameters to allow the client and server to agree on a premaster secret

 Exchange certificates and cryptographic information to allow the client and server to authenticate themselves

– TLS 1.2 handshake

Generate a master secret from the premaster secret and exchanged random values

- Provide security parameters to the record layer
- Allow the client and server to verify that their peer has calculated the same security parameters and that the handshake occurred without tampering by an attacker

– TLS 1.2 handshake

 *Session resumption is useful because the handshake is so heavy (back and forth, cipher selection, random value generation, encryption etc.

 Instead of recalculating every value, if a client has a session ID, the server and client can skip to ChangeCipherSpec and start sending encrypted traffic

– TLS 1.2 handshake

function.k(a,c,d)[if]

Ignore keyless part

Note both parties have their own private DH secret

SSL Handshake (Diffie-Hellman) Without Keyless SSL





– Enter Diffie-Hellman algorithm





https://en.wikipedia.org/wiki/Diffie%E2%80%93Hellman_key_exchange

– Diffie-Hellman key exchange

– <u>Logjam vulnerability</u>

- <u>See also cloudflare</u>



- TLS 1.2 handshake - Illustrated Guide

For examples see the following wireshark captures:

 sudo puppet agent --test --no-daemonize --server instructor.itec2210.ca

tshark -r /home/itec2210/tls.pcap

Trace through and see each ClientHello with different cipher suites

– TLS 1.2 handshake

Note you'll see many ECDHE, DHE, ECDH cipher suites

 EC means elliptic curve, which is a branch of cryptography that uses curves and functions to select primes

 Much smaller primes needed when using curves, means faster TLS negotiation, fewer bytes in a handshake

– TLS 1.2 handshake

Something like DHE/EDH means Ephemeral Diffie-Hellman

 No two TLS sessions will share the same key, meaning if a session key is captured, no others are compromised

Principle here is called <u>Perfect Forward Secrecy</u>

<u>https://scotthelme.co.uk/perfect-forward-secrecy/</u>